

# **A Versatile Hardware Architecture for a Constant False Alarm Rate Processor Based on a Linear Insertion Sorter**

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## **Abstract**

Constant False Alarm Rate (CFAR) algorithms are used in digital signal processing applications to extract targets from background in noisy environments. Some examples of applications are target detection in radar environments, image processing, medical engineering, power quality analysis, features detection in satellite images, Pseudo-Noise (PN) code detectors, among others. This paper presents a versatile hardware architecture that implements six variants of the CFAR algorithm based on linear and non-linear operations for radar applications. Since some implemented CFAR algorithms require sorting the input samples, a linear sorter based on a First In First Out (FIFO) schema is used. The proposed architecture, known as CFAR processor, can be used as a specialized module or co-processor for Software Defined Radar (SDR) applications. The results of implementing the CFAR processor on a Field Programmable Gate Array (FPGA) are presented and discussed.

**Keywords:** CFAR Processor, Software Defined Radar, Hardware Architecture

# 1 Introduction

The problem of detecting target signals in background noise of unknown statistics is a common one in sensor systems such as radars and sonars. In radar applications, this noise usually comes from thermal noise (receiver's noise), clutter, pulse jamming or other undesired echo received at the antenna. Adaptive digital signal processing techniques are often used to remove noise and to enhance the detectability of targets in many situations. It is important for radar processing systems to operate in non-stationary background noise environments with a predetermined constant level of performance, i.e., in signal processing terms; the objective is to maintain a constant false alarm rate. A solution to overcome the problem of noise added to the target signal is to use the constant false alarm rate (CFAR) algorithm, which sets a threshold adaptively, based on local information of total noise power. The threshold set by the CFAR algorithm is obtained on a sample by sample basis, using estimated noise power by processing a group of samples surrounding the sample under investigation [1] and [2].

Several variants of the CFAR algorithm have been proposed in the radar literature to deal with different problems present in radar applications. These techniques require linear operations (such as getting the maximum, minimum, or average of a set of values) or nonlinear operations like sorting a set of values and selecting one on a specific position before performing a linear operation. These different techniques have been developed in order to increase the target detection probability under several environment conditions, especially to deal with two of them: regions of clutter

transitions and multiple target situations. The first situation occurs when the total received noise power within a single reference window changes abruptly, leading to excessive false alarm or target masking. The second situation is encountered when there are two or more closely spaced targets in the reference cells, leading the CFAR processor to report only the strongest of the targets, i.e., there is a target masking [2].

Although theoretical aspects of detection in radar systems, including the CFAR algorithm, are very advanced [1][2][6], and analog implementations have been used in radar systems for a several years, recent technology developments have made it practical to explore digital implementations of CFAR and other algorithms to support the SDR paradigm. SDR systems can be implemented using digital systems to accommodate various radar sensors for different detection conditions. This means that they can be designed with the capability for being configured at run time by means of selecting an appropriate operation mode according to current environment conditions [10] [11].

Real-time performance of radar detection algorithms as required by modern radar systems can be achieved by state-of-the-art general purpose processors or digital signal processors (DSPs) optimized to deal with high computational load and high input/output data rates. As an alternative, these requirements can be met by specialized custom architectures that exploit the parallelism on the algorithm employed [12]. For practical SDR applications all processing blocks, including CFAR, must support several processing modes and operate with a high computational load in real-time. This work

presents a versatile custom hardware architecture that supports six variants of the CFAR algorithm and is suitable for SDR applications. The proposed CFAR hardware architecture, or CFAR processor, can be used as a specialized processing module or co-processor in the receiver's processing chain of a SDR system. The supported processors differ in the method they use to obtain the adaptive threshold from the reference window. Three of the implemented processors are based on linear operations while the other three are based on nonlinear operations that in turn are based on rank order statistics. This rank order operation means sorting and then selecting a  $k$ -th value from a set of values, which adds computational complexity to the CFAR processors. Because of this, a linear insertion sorter schema is used to keep the data sorted and then to perform the desired operation.

## **2 Detection Overview**

A radar transmitter generates an electromagnetic signal that is broadcast to the environment by an antenna. A portion of the energy of this broadcast signal is reflected by targets. This reflected energy is received by the same antenna and sent to the receptor. At the receptor this energy is digitalized to produce raw data that is then processed to obtain the desired target information. Figure 1 shows a radar receiver processing chain and the position of the CFAR processor.

**Figure 1. Traditional signal processing radar chain.**

The radar detection of echo signals from targets is performed by establishing a threshold at the output of the receiver. If the receiver output is greater than the established threshold, it is declared a *target presence*; otherwise it is declared a *target absence* i.e., only noise is present. This noise presence can be caused by weather conditions (clutter), thermal noise from radar devices, pulse jamming or interference.

If the fixed threshold level is set properly, the receiver output would not exceed the threshold if only noise is present, but the receiver output would exceed this threshold if along with noise, a target is present. If the threshold level were fixed too low, noise alone might exceed it, situation called *false alarm*. If the threshold is fixed too high, only strong target echoes would be able to exceed it and weak target echoes might be not detected, a situation called *missed detection*. In early radars, this threshold level was set based on the radar operator judgment. In figure 2 there is an example of the output of a radar receiver as a function of time, which is called *range profile*. The fixed threshold line is represented by the horizontal dashed line and the adaptive threshold is represented by the dashed-dotted line. For this example, suppose that a signal has four segments A, B, C and D. The signal segments A and B are targets plus noise and C and D are noise alone. The signal segment A is detected correctly by this fixed threshold, while B is not strong enough to be detected and therefore it is missed. The noise from signal segment B is weak because of the negative noise that was added to the original signal strength. Signals segments C and D are false alarms, and they are increased because of the presence of positive noise. B could be detected if the fixed threshold was lower, but this might increase the false alarms. The selection of a proper fixed threshold

is a compromise that depends upon how important it is to avoid the mistake of failing to declare a target presence (*missed detection*) or falsely indicating the presence of a target when none exists (*false alarm*) [1]. On the other hand, with an adaptive threshold, the signal segments A and B, which exceed the threshold level, can be detected correctly, while C and D are assumed to be noise because they do not exceed the threshold.

**Figure 2. Example of a radar receiver range profile with fixed and adaptive threshold.**

In order to maintain the false alarm rate at a constant value, the threshold has to be varied adaptively. This is achieved when the CFAR processor automatically raises the threshold level, thus avoiding overload of the automatic detection system. A constant false alarm rate is achieved at the expense of a lower detection probability of desired targets. Also, CFAR produces false echoes when there is nonuniform clutter, suppresses nearby targets and decreases the range resolution. The detection probability and false alarm probability are specified by the system requirements.

### **3 CFAR and OS-CFAR processors**

The need for CFAR processors was recognized when the early automatic detection and tracking systems were installed as add-ons to existing radar with not moving target indicators (MTI) or relatively poor MTI that did not have a good clutter rejection. CFAR is needed for maintaining operation for automatic detection and tracking systems. If CFAR were not used in radar, it would cause an excessive number of false

alarms, due to noise or clutter, degrading the automatic detection and tracking system performance.

### **3.1 Algorithms description**

Several algorithms have been proposed for building CFAR processors. The most commonly used CFAR processors are the cell averaging (CA-CFAR) processor proposed in [7], the greatest of (GO-CFAR) processor proposed in [8], the smallest of (SO-CFAR) processor proposed in [9], the generalized order statistics cell averaging (GOSCA-CFAR), generalized order statistics greatest of (GOSGO-CFAR), and generalized order statistics smallest of (GOSSO-CFAR) processors proposed in [5]. Figure 3 shows a general block diagram of a generic CFAR processor.

#### **Figure 3. Generic CFAR processor.**

This processor consists of a reference window with  $2n$  cells that surround the cell under test. Each cell stores an input sample and such values are right shifted when a new sample arrives. Some  $m$  guard cells are incorporated in order to avoid interference problems in the noise estimation. The spacing between reference cells is equal to the radar range resolution (usually the pulse width). The reference cells are used to compute the  $Z$  statistic and, depending on the technique, this operation can be linear or nonlinear. The  $Z$  statistic and a scaling factor  $\alpha$  are used to obtain the threshold. This scaling factor depends on the estimation method applied and the false alarm required according to the application. It is also related to the noise distribution in the radar environment. The resulting product  $\alpha Z$  is used as the threshold value that is compared with the cell under

test (CUT), to determine whether the CUT is declared a target. The target detection problem can be modeled by the following hypothesis testing problem:

$$\begin{aligned} H_1 : y &= d+g \\ H_0 : y &= g \end{aligned} \tag{1}$$

where  $H_1$  and  $H_0$  are the target present and target absent hypothesis, respectively;  $d$  represents the target signal and  $g$  the environmental noise component. The decision criterion is represented by:

$$e(y) = \begin{cases} H_1, & CUT \geq \alpha Z \\ H_0, & CUT < \alpha Z \end{cases} \tag{2}$$

If the values of the CUT exceed the  $Z$  statistic, then a target presence is declared, *i.e.* the CFAR processor outputs 1 if a target is present, otherwise it outputs 0:

$$e(y) = \begin{cases} 1, & CUT \geq \alpha Z \\ 0, & CUT < \alpha Z \end{cases} \tag{3}$$

The method to obtain the  $Z$  statistic from the reference window might be based on linear or nonlinear operations. The most common linear processors are the CA-CFAR, GO-CFAR and SO-CFAR. These processors calculate the arithmetic mean of the amplitude contained in the  $Y_1$  lagging cells and  $Y_2$  leading cells from the CUT. The CA processor estimates the arithmetic mean, the GO and SO take the major and minor values of  $Y_1$



and  $Y_2$ , respectively. Equation 4 summarizes these three linear operations for the  $Z$  statistic.

$$e(y) = \begin{cases} \frac{1}{2} (Y_1 + Y_2) \\ \max (Y_1, Y_2) \\ \min (Y_1, Y_2) \end{cases} \quad (4)$$

Among the nonlinear processors are the OSCA-CFAR, OSGO-CFAR and OSSO-CFAR [4]; and their generalized form called GOSCA-CFAR, GOSGO-CFAR and GOSSO-CFAR processors [5]. These order statistics processors need to perform a rank-order operation over the leading and lagging reference cells, i.e., sort the reference cells values and then select the  $k$ -th sorted value. The rank-order parameter  $k$  can be deliberately selected among the sorted values. The GOSCA-CFAR, GOSGO-CFAR and GOSSO-CFAR processors, perform the selection of the  $k$ -th ( $Y_{(1)}$ ) and  $i$ -th ( $Y_{(2)}$ ) sorted value from the leading and lagging cells, respectively. Once these two values have been selected, the  $Z$  statistic is calculated in a similar way to the linear processors as shown in the following equation:

$$e(y) = \begin{cases} \frac{1}{2} (Y_{(1)} + Y_{(2)}) \\ \max (Y_{(1)}, Y_{(2)}) \\ \min (Y_{(1)}, Y_{(2)}) \end{cases} \quad (5)$$

The difference between these generalized processors and OSCA-CFAR, OSGO-CFAR and OSSO-CFAR processors is that the last ones only perform the selection of the  $k$ -th sorted value from both the leading and lagging cells. Therefore, OSCA-CFAR, OSGO-CFAR and OSSO-CFAR can be considered a special case of their generalized counterpart when  $k = i$ .

### **3.2 Features of CFAR algorithms**

CA-CFAR processor performs better than other CFAR processors, in terms of detection probability, when operating in a homogeneous background when reference cells contain independent and identically distributed observations governed by an exponential distribution. However, the CA-CFAR detection performance degrades in multiple target situations and regions of power transitions. Both situations result in an excessive number of false alarms i.e. an inferior behavior in nonhomogeneous situations [2]. GO-CFAR and SO-CFAR processors came up in order to solve these problems. GO-CFAR maintains a constant false alarm rate at clutter edge and regions of power transitions; meanwhile, SO-CFAR resolves the primary target in multiple target situations when all the interferers are located in either the leading or lagging cells. However, GO-CFAR detection performance in multiple target situations is poor and SO-CFAR has undesired effects when interfering targets are located in both halves of the reference cells [2].

To improve the detection performance on these situations, order statistics techniques (OS-CFAR) were proposed in [3] reporting better overall performance results. In [4], two modified OS-CFAR processors that require less processing time than the OS-CFAR

processor were proposed. The GOSCA-CFAR, GOSGO-CFAR and GOSSO-CFAR processors are presented in [5]. These processors achieve better detection performance in the presence or in the absence of interference. The generalized processors are more robust than the processors proposed in [3] and [4] because of the selection of the  $k$ -th and  $i$ -th rank-order sample and, because they require half of the time for sorting compared with OS-CFAR. Among these generalized processors, GOSCA-CFAR processor obtains the best detection performance in both homogeneous background and multiple target situations. The GOSGO-CFAR has the same performance as the OS-CFAR, and provides good transition clutter protection. If the number of interfering targets in the reference cells is equal to the greatest interfering targets allowed, the performances of these generalized processors is better than the OS-CFAR processor.

Each one of the explained processors has advantages and drawbacks, and may be optimal under particular environment conditions. The detection performance is altered by varying the number of references cells, guard cells, the CFAR processor, the  $k$ -th rank-order sample, and the false alarm required; represented by the scaling factor  $\alpha$  [2]. In order to give robustness to the target detection process in radar systems, a specialized architecture that supports several of these processors and allows changing their operating parameters such as: CFAR algorithms, scaling factor  $\alpha$ , and the  $k$ -th and  $i$ -th rank-order sample would be highly beneficial. The following subsection reviews some related works found in the literature.

### **3.3 CFAR processors review**

In radar related literature; there are few reported hardware architecture implementations of CFAR processors. An architecture for CA-CFAR, GO-CFAR and SO-CFAR processors is presented in [12]. This architecture implements the average computations with two accumulating processing elements (APE) and a configurable threshold processing element (CTPE). These APEs compute the sum of their corresponding reference cells (lagging and leading) and the CTPE computes the threshold operation (cell-averaging, greatest of and smallest of). This architecture uses 12 bits for data, 32 reference cells, 8 guard cells and the internal temporal data of 18 bits precision in the accumulator for the worst case, and its operating frequency is 120 MHz on a XC2V250 Virtex II FPGA device. In [13], authors presented an FPGA-based implementation of a matched filter with an OS-CFAR processor, focused on the adaptive pseudo noise (PN) code acquisition. In this work, the OS-CFAR processor uses the bubble sorting algorithm to find the  $k$ -th biggest sample in the reference cell. This OS-CFAR processor uses 16 bits for data, 16 reference cells and none guard cell. The resulting architecture is implemented on a XCV400E Virtex E FPGA device with a maximum clock frequency of 205 MHz. A specialized architecture of a CA-CFAR processor is presented in [14]. This architecture is implemented on a XC9600 FPGA device and it was tested with 8 bit for each of the 16 reference cells. Its implementation consists of a storage circuit, 17 8-bit shift registers, two accumulator circuits, eight 8-bit adders, a multiplier circuit and one 8-bit comparator. Area results and frequency operation are not reported.

A systolic architecture for CFAR processors was presented in [15]. In this work, authors proposed a systolic array capable of supporting two OS-CFAR processors: CMLD

(Censored Mean Level Detector) and MXCMLD (Max Censored Mean Level Detector) CFAR processor. The first processor censors the largest reference samples sorted and then adds the remaining samples; meanwhile the second one gets the maximum of the sum of the remaining samples from the lagging and leading cells, after being censored. A modified OS-CFAR systolic architecture is presented in [16]. This architecture requires fewer PEs, interconnections, and gates, and reaches twice the throughput when compared against the architecture in [15]. In [16], authors also proposed a systolic architecture for the OSGO-CFAR and OSSO-CFAR processor by slightly modifying the OS-CFAR systolic architecture. The throughput of this architecture is the same as the original proposed OS-CFAR. In [17], a parallel systolic architecture for a CFAR processor with adaptive post-detection integration (API) is presented. The CFAR processor for this architecture was developed, analyzed and synthesized in a systolic architecture in order to detect targets in presence of pulse jamming. This architecture has a linear structure, specially designed for real-time implementation. It uses four sequential blocks that perform the operations of sorting, censoring, integration and comparison. These four blocks are constructed by five processing elements, which perform different logical operations.

A patented system called ES-CFAR (Expert System) is presented in [18]. This system senses the clutter environment and by means of a set of rules along with a voting scheme, selects the most appropriate CFAR processor(s) to produce detection decisions that will outperform a single processor. This expert system is based on five separate CFAR processors: CA-CFAR, GO-CFAR, SO-CFAR, OS-CFAR and TM-CFAR

(Trimmed Mean). The ES-CFAR system is based on the use of knowledge-based signal processing algorithms [19].

## **4 Proposed CFAR Processor Architecture**

The proposed architecture uses a linear sorter that performs the rank-ordering operation needed in the order statistic processor. Since keeping the values sorted does not affect the averaging process needed in the linear processor, the use of a linear sorter is possible. The architecture is parameterizable in terms of the number of reference and guard cells, and arithmetic precision.

### **4.1 Linear Insertion Sorter**

The linear sorter used in this architecture implements the insert sort algorithm presented in [20]. It consists of an array of identical processing elements (PE) which are called Sorting Basic Cell (SBC). In order to fulfill the FIFO sorting functionality, the SBCs must be interconnected in a simple linear structure, called sorting array. This array is shown in figure 4, where  $R_n$  represented the  $n$ -th SBC inside the sorting array and  $D$  is the incoming value.

#### **Figure 4. Sorting array.**

This SBC array sorts the values as they are introduced into the sorting array, discarding the oldest value, while maintaining the values sorted in a single clock cycle i.e. in a

FIFO fashion. The SBC has a register with synchronous load to store the value ( $X_i$ ), a counter with synchronous reset and load to store the life period of the value ( $CNT_i$ ), a comparator, four 2-input-1-output multiplexers and control logic (figure 5).

### Figure 5. Architecture of the SBC.

The SBC control logic consists of four Boolean equations, which control the register, the counter and the multiplexers:

$$load = (p_i \oplus cnt_{i+1}) + expired_i \quad (6)$$

$$LR = p_i \bullet load \quad (7)$$

$$reset = load \bullet [(p_{i-1} \bullet \bar{p}_i) + (p_i \bullet \bar{p}_{i+1})] \quad (8)$$

$$cnt_i = cnt_{i+1} + expired_i \quad (9)$$

where  $p_i$  is the comparator output. The signals  $p_{i+1}$  and  $p_{i-1}$  correspond to the right and left SBC neighbors respectively and  $cnt_{i+1}$  is the flag coming from the SBC immediately to the right. This signal helps to detect if the life period value of one SBC to the right has expired. The signal  $expired_i$  indicates when the life period value has expired inside of one SBC and only one of the  $expired_i$  signals in the sorting array is activated in a clock cycle. In order to perform the insert sort algorithm correctly, the left most  $p_{i+1}$  signal's value is always 1 and the right most  $p_{i-1}$  signal's value is 0. This can be viewed as the left most value having the largest value while the right most has the smallest one. To ensure proper behavior, all  $R_i$  registers must be initialized to zero, while life counter

values CNT must be initialized according to  $CNT_i = i$ , i.e. its position within the sorting array. This linear insertion sorter architecture is fully described in [21].

## 4.2 CFAR Processor Architecture

The proposed CFAR processor architecture, figure 6, has two SBC's sorting arrays for  $2n$  reference cells,  $2m+1$  shift registers for the guard cells and CUT, which is at the middle of these registers. Also, the architecture has two  $n$ -input-1-output multiplexers that perform the rank operation for the lagging and leading sorting arrays. Given that the reference cells values are ordered, the  $k$ -th and  $i$ -th value can be selected by the control signals Sel-k and Sel-i respectively. The result of this selection is the  $Y_{(1)}$  and  $Y_{(2)}$  values needed in the nonlinear operations shown in equation 5. Considering that the function performed by each of the  $2n$  SBC. The number of arithmetic operations performed by the architecture per clock cycle is  $2n+7$ . This number is obtained considering that each SBC performs a comparison operation ( $2n$ ) and each accumulator performs one add and one subtract operation ( $2+2$ ). The other three operations are: 1) one ALU operation, 2) the product  $\alpha Z$  and, 3) the threshold comparison to obtain  $e(y)$ .

### Figure 6. CFAR Processor hardware architecture.

For the linear operations presented in equation 4, it is required to add all values stored in the leading and lagging sorting arrays for computing the average. In order to perform this operation, it is not necessary to add all values each time that one value from the



sorting array is inserted and deleted. Once a value is inserted and other one deleted, the preceding result can be used to compute the next result without adding all values. Only by adding and subtracting the newest and oldest values respectively, the next result is obtained. This whole operation can be performed by the PE Accumulator, which computes the average of  $Y_1$  and  $Y_2$  values on each sorting array. The PE Accumulator, figure 7, consists of an adder, which receives the incoming value, a subtracter, connected to one of the multiplexers which select the oldest value stored in the sorting array, a register to store the accumulated value and a left shifter that performs the division needed to compute the average  $Y_n$  value. Because of the left shifter, the number of reference cells must be a power of two. This does not restrict the usability of this architecture as the number of reference cells used in practical applications is usually a power of two [1].

### **Figure 7. Processing element accumulator.**

The oldest value stored in the sorting array is obtained by a  $n-1$  multiplexer whose control line value is generated by a priority decoder. The input of this decoder is a data bus formed by the  $expired_n$  signals coming from the SBC in the sorting array and the output bus is *SelOldest* (Select Oldest), as shown in figure 8. The *expired* signal indicates when the life period value has expired in only one SBC's, i.e., the oldest value that must be subtracted and passed to the shift registers.

### **Figure 8. Priority decoder and its connection with the sorting array.**

Two 2-1 multiplexers perform the selection between  $Y_1$  and  $Y_{(1)}$  from the lagging window and  $Y_2$  and  $Y_{(2)}$  from the leading window. The selection is performed by the one bit signal *SelOp* (Select Operation). An ALU-like module provides the three modalities for computing the  $Z$  statistic: the average, the maximum and the minimum of the rank-order or the accumulated value. The desired modality is chosen using the control signal *SelDet* (Select Detector) established either manually by the user or by an automatic control expert system. A multiplier scales up the  $Z$  statistic with the scaling factor  $\alpha$  and a comparator decides whether a target is present or absent as indicated by equation 3.

The control signals, *SelDet* and *SelOp* can be grouped into a data bus in order to perform the selection among the six detectors with their corresponding operations with only a data bus. The arithmetic precision of scaling factor  $\alpha$  can be modified in off-line work.

In this architecture, at each clock cycle, values flow from the lagging sorting array to the shift registers and to the leading sorting array. In order to begin the target detection processing, a reset signal must be applied to the SBC in order to initialize the counters inside of the SBCs and PE Accumulators. Once the data begins to flow,  $2n+2m+1$  clock cycles are required for having all the values stored in the sorting arrays and in the shift registers. After this latency time, the architecture produces a valid output each clock cycle, allowing the continuous operation of the target detection process.

## 5 Results and Discussion

The architecture was developed for a commercial X-band non-coherent radar. The radar performs a 360 degrees scan in 2.5 seconds (24 rotations per minute). Incoming raw data from the receiver are sampled to produce a set of 4096x4096 samples per scan that are processed in a stream fashion to generate the radar image. For the purpose of validation, the proposed architecture was modeled using the VHDL Hardware Description Language and synthesized with Xilinx ISE 9.1 targeted for a XtremeDSP Development Kit, that includes a Xilinx's Virtex-4 XC4VSX35 FPGA device. It is important to emphasize that the architecture was implemented in an FPGA device for the purpose of validation. This allowed us to obtain simulation and performance results in order to show that this compact custom architecture is able to support several variants of the CFAR algorithm while maintaining high performance for demanding radar applications. As no FPGA exclusive resources or features were used to design the architecture, it can be implemented as a custom architecture or used as a coprocessor attached to a general purpose processor.

## 5.1 Area and Performance Results

The architecture was synthesized for several configurations in terms of the number of reference and guard cells. For throughput comparison with the architecture presented in [12] that reports a throughput of 840 MOPS, the proposed architecture was synthesized for a Virtex II FPGA device. The throughput achieved using 32 reference and 4 guard cells is 7,526 MOPS which is nine times better than the throughput reported in [12] using the same number of cells. In order to show results for a more up to date device, the proposed architecture was also synthesized for a Xilinx's Virtex-4 XC4VSX35 FPGA device using different number of reference cells. Table 1 summarizes the results in terms of FPGA hardware resources utilization including four guard cells at each side of the CUT. Table 1 also shows the throughput measured in millions of operations per second (MOPS) and the required processing time for the radar data set of 4096x4096 samples. All these four configurations use 12-bit to represent input data. The third row presents the results obtained for the most common configuration for radar applications [1]. Using 32 reference cells, the proposed architecture requires 84 milliseconds to process a radar data set of 4096x4096 samples, which is 30x times faster than the required theoretical processing time of 2.5 seconds needed for this application parameters; thus this module can be potentially used in radars with much higher resolution or CFAR processor with larger n and m values, i.e., the amount of reference and guard cells. In fact, with a greater configuration of the CFAR detector of 12-bits of data, 64 reference cells and 8 guard cells, the architecture has a maximum operation

frequency of 165 MHz, requiring 99 milliseconds to process the same data amount. This configuration is 25x times faster than the theoretical required processing time.

**Table 1. Resources utilization for a Xilinx’s Virtex-4 XC4VSX35 FPGA device**

The throughput in Table 1 was calculated considering that the architecture performs  $2n+7$  arithmetic operations per clock cycle as described in subsection 4.2. Considering the design with 32 reference cells at the maximum operation frequency, this architecture achieves a throughput of 14,058 (MOPS). This is the result of the sorting schema used to support the rank-order operation, since each SBC used for storing the values in a sorted way requires two operations and the number of SBC used is proportional to the number of reference cells used.

For performance comparison, consider the results presented in [12], where a CA-CFAR processor was implemented on a TMS320C6203 DSP device. The processing time obtained for this DSP implementation was 420 milliseconds to process a sample data set of the same size. It is important to mention that the implementation in [12] does not perform the sorting operation.

A fair comparison with other reported CFAR processors in terms of area and performance is not feasible because they do not perform the same functionality as the proposed architecture. However, in order to give an idea of the required hardware resources, Table 2 shows a comparison of the proposed CFAR hardware processor against other works in terms of the number of hardware elements they require. In this

table  $p$  refers to the amount of cells in a reference window, i.e., the sum of reference cells ( $2n$ ), guard cells ( $2m$ ) and CUT,  $p = 2n+2m+1$ . The clock cycles refer to the latency required to process the data and the delay elements refer to hardware elements that are needed to store data. The architectures described in [15], [16] and [17] are based on systolic designs. In [17], the architecture implements a two dimension CFAR processor ( $p$  rows and  $q$  columns). For the purpose of generating Table 2 it was assumed that  $q = 1$ .

### **Table 2. Comparison with others CFAR Processor architectures.**

The architectures in [15] and [16] implement the OS-CFAR, OSGO-CFAR and OSSO-CFAR processors. In [17], a censored technique over the sorted samples is used and then the uncensored samples are added. The architecture presented in [12] implements the CA-CFAR, GO-CFAR and SO-CFAR processors in a specialized architecture. Even though the proposed architecture implements a sorter functionality, it has the same latency than the architecture in [12] which implements three of the six processors that the proposed architectures does. Also, both architectures use the same amount of adders and multipliers. Due to the proposed CFAR architecture performs a sorting functionality it uses more comparators than [12], which does not perform this functionality. Among the other three systolic architectures, the proposed CFAR architecture uses the least number of comparators, delay elements and adders, but it uses more multiplexors. However, when compared with the systolic architectures presented

in [15] and [16] the proposed solution implements the OSGO-CFAR and OSSO-CFAR processors with only a slightly higher amount of hardware elements.

## 5.2 Architecture Validation

To validate the results of the proposed CFAR processor, the six selected variants of the CFAR algorithm were modeled in software using C language. Input data, i.e. range profiles, were obtained from real radar scans. Each radar scan consists of several range profiles. Figure 9 shows a range profile example with 500 samples. Both software and hardware models were fed by the radar data and their corresponding outputs were compared, i.e. when the CFAR hardware architecture performs an erroneous detection or miss detection.

### Figure 9. Radar receiver range profile example.

For these tests a typical CFAR configuration was used, i.e. 12-bits for data, 32 reference cells and 8 guard cells and the  $k$ -th and  $i$ -th rank-order samples = 12. According to [2] a value of  $0.75n$  has the best detection performance, where  $n$  is the size of the lagging and leading windows. The value used for the scaling factor was  $\alpha = 0.95$ . However, an exact fixed point representation for the scaling factor is not suitable for the CFAR hardware detector. Therefore, a scaling factor of  $\alpha = 0.9501953125$  was used in the proposed architecture. This approximation is very close to the value of 0.95 which was used for the software implementation.

Figure 10 shows the resulting threshold calculated by the linear CFAR detectors implemented in software (figure 10a) (C language) and hardware (figure 10b), using as input the radar receiver range profile shown in figure 9. At first glance, there are not significant differences among the threshold calculated by the different implementations. However, there are small differences between the thresholds calculated in software and hardware. The differences at each sample between software implementation and the proposed CFAR hardware architecture are shown in figure 10c. Note that the x-axis scale, indicating the difference between the software implementation and the CFAR hardware architecture, is in mV meanwhile the thresholds calculated for both implementations are in V. This difference is caused by an error in the scaling factor fixed point representation and the average computing need on the linear detectors. A similar figure can be viewed in figure 11, where resulting thresholds calculated by the nonlinear CFAR detectors implemented in software (figure 11a), hardware (figure 11b) and its corresponding differences (figure 11c). Results show that the error in the calculated threshold is in the order of  $10^{-3}$  which is not significant considering that these algorithms are based on probability of detection [1].

**Figure 10. Radar receiver range profile and thresholds calculated by the linear CFAR processor implemented in (a) software and (b) hardware. The differences among each one of the CFAR detectors implementations in shown in (c).**

**Figure 11. Radar receiver range profile and thresholds calculated by the nonlinear CFAR detector implemented in (a) software and (b)**



**hardware. The differences among each one of the CFAR detectors implementations in shown in (c).**

Figure 12 illustrates how a threshold calculated by the proposed CFAR hardware architecture varies when switching between different CFAR algorithms. This feature is needed in a SDR system which needs to change its parameters in run time execution. For clarity purposes, figure 12 only shows the thresholds without the radar echoes as shown in previous figures. Each line style represents a different detector. Figure 12a shows the threshold calculated during 500 samples, from sample 1 until sample 500. Note that the first 250 samples are calculated by OSGO-CFAR algorithm (crossed line) and the 250 remaining samples are calculated by GO-CFAR algorithm (pointed line).

**Figure 12. Thresholds obtained by switching between CFAR algorithms.**

Figure 12b illustrates the threshold obtained in two runs of the simulation. The first run corresponds to the one shown in Figure 12a, i.e. using OSGO-CFAR for the first 250 samples (crossed line) and then GO-CFAR (pointed line). For the second run only GO-CFAR was used. For the first 250 samples the run using GO-CFAR (pointed line) differs from that using OSGO-CFAR (crossed line). Note after sample 250, both runs are identical. This shows how a switch between algorithms can be done in a single clock cycle without a transition phase. This is possible because the architecture always calculates the values needed to obtain the threshold for all variants of the CFAR

algorithm, and switching algorithms implies only to select the appropriate values to obtain the next threshold value.

## **6 Conclusions**

CFAR processors are used in signal processing applications to extract target signals from noisy background. For radar applications, a number of CFAR algorithms have been proposed for several environment conditions. This paper describes a versatile processing architecture that allows switching among six CFAR algorithms and operating parameters in a single clock cycle in order to provide robustness to the target detection process. This work has shown that it is feasible to implement a single architecture to support several variants of the CFAR algorithm while achieving the data rates and algorithmic performance required by demanding radar applications. This work has also shown that modern FPGAs devices are well suited to efficiently implement the target detection process of conventional and software defined radar systems.

## **7 Acknowledgments**

First author thanks the National Council for Science and Technology from Mexico (CONACyT) for financial support through the scholarship number 204500.

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<b>Ref. Cells Amount</b>	<b>Slices Count</b>	<b>LUTs Count</b>	<b>FF Count</b>	<b>Speed (MHz)</b>	<b>Throughput (MOPS)</b>	<b>Processing Time (mS)</b>
<b>8</b>	315	266	581	256	5,888	64
<b>16</b>	602	423	1,147	218	8,502	77
<b>32</b>	1,364	690	2,637	198	14,058	84
<b>64</b>	2,790	1,260	5,430	165	22,275	99

**Table 1. Resources utilization for a Xilinx's Virtex-4 XC4VSX35 FPGA device**

Number of	CFAR Processor Hardware Complexity				
	Hwang [15]	Han [16]	Behar [17]	Torres [12]	Proposed Architecture
Comparators	$2p+1$	$2p+1$	$3(p^2+p+4)/2$	2	$2n+2$
Multiplexors	$2p+1$	$2p+1$	$p^2+p+2$	1	$8n+7$
Delay Elements	$8p+5$	$5p+7$	$2p^2+2p+1$	$p$	$4n+2m+3$
3-input Adders	0	$p-1$	0	0	0
2-input Adders	$2p$	2	$p+2$	3	3
Multiplicator	1	1	1	1	1
Clock Cycles	$2p+3$	$p+3$	$2p+2$	$p$	$p$

**Table 2. Comparison with others CFAR Processor architectures.**

**Figure 1**  
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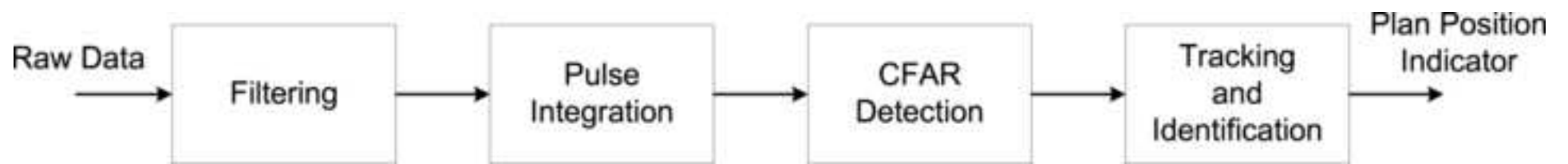




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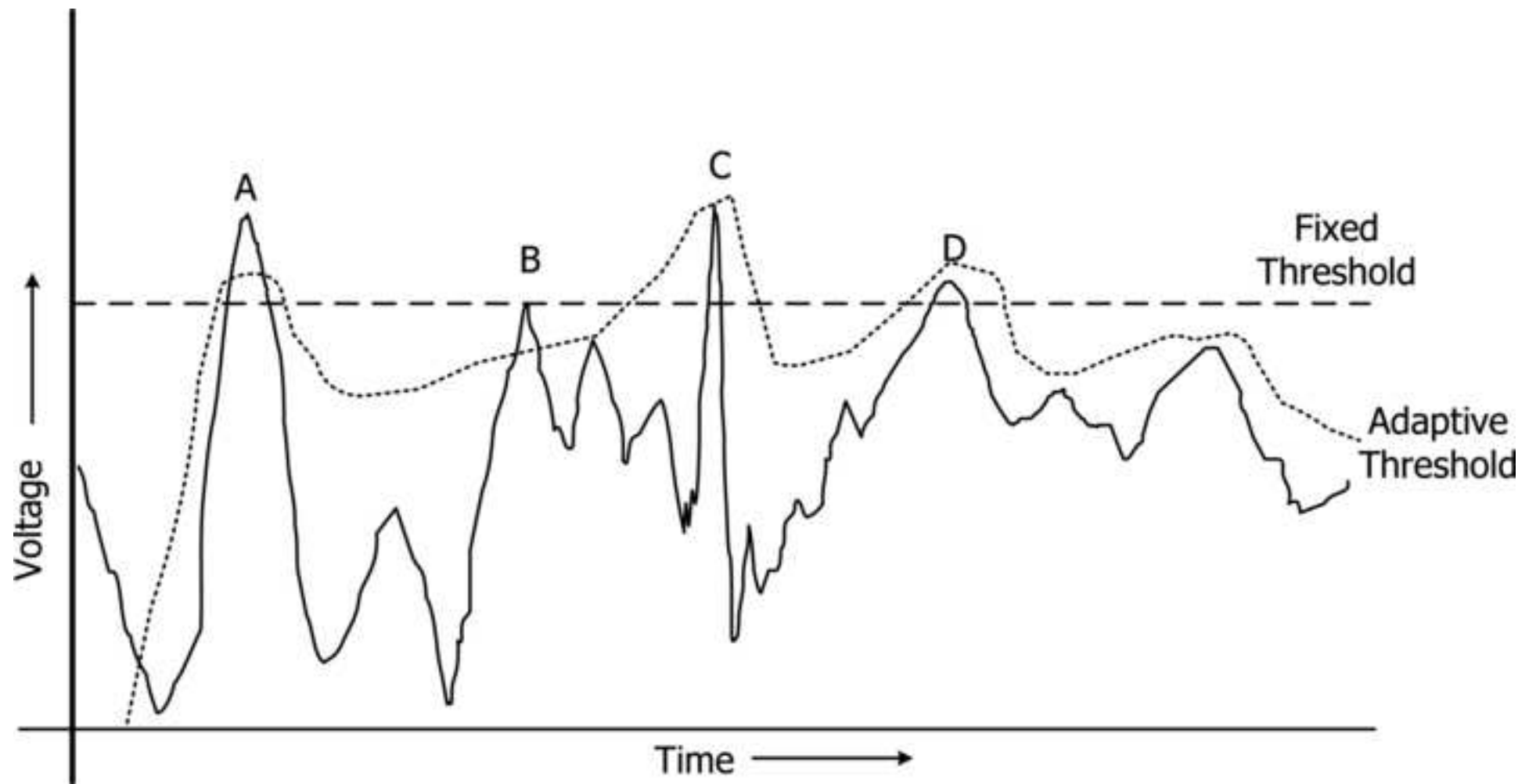


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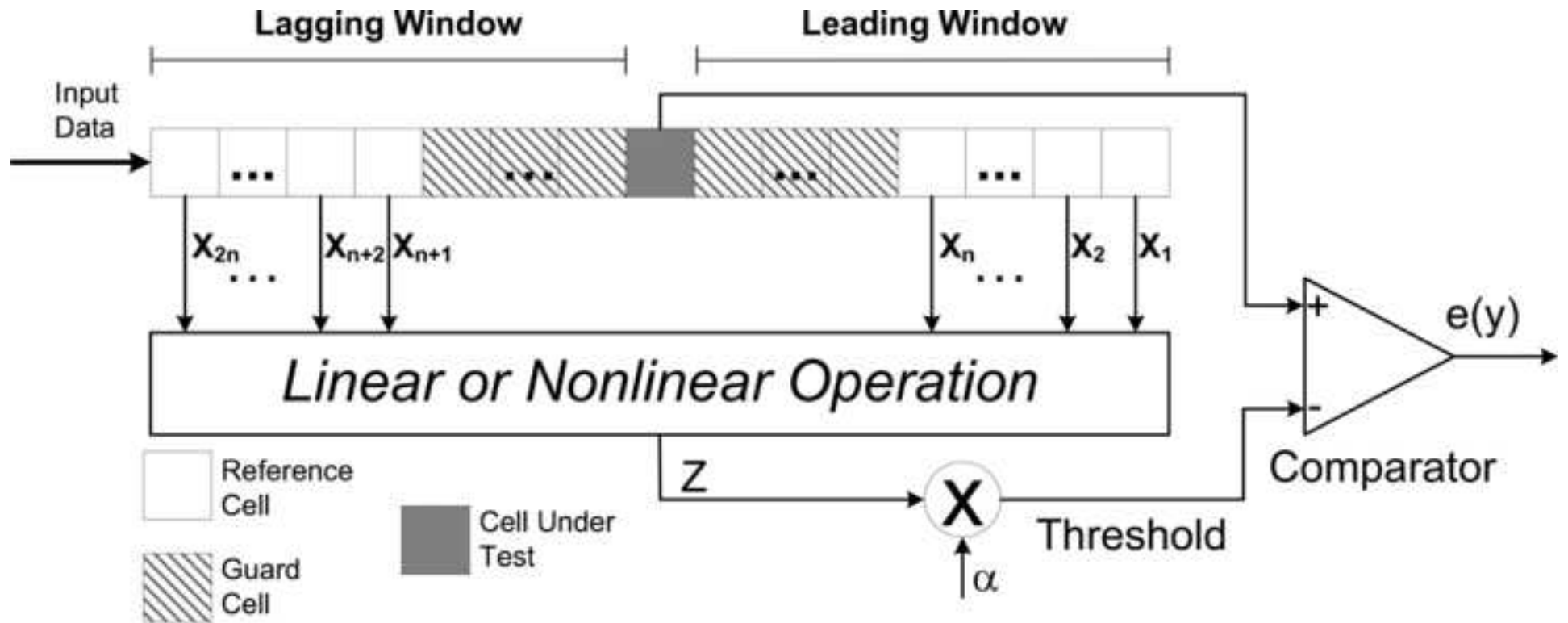


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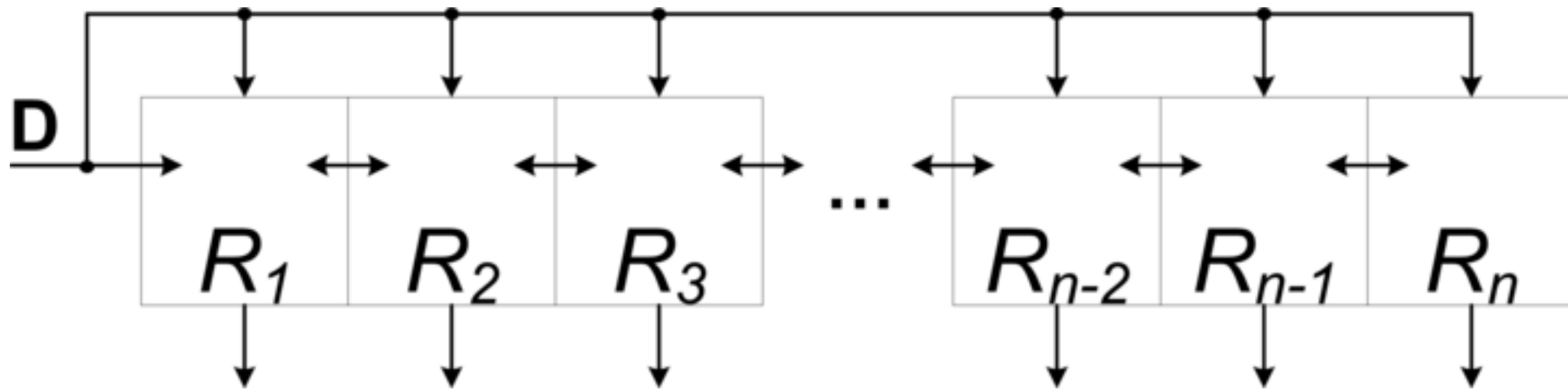


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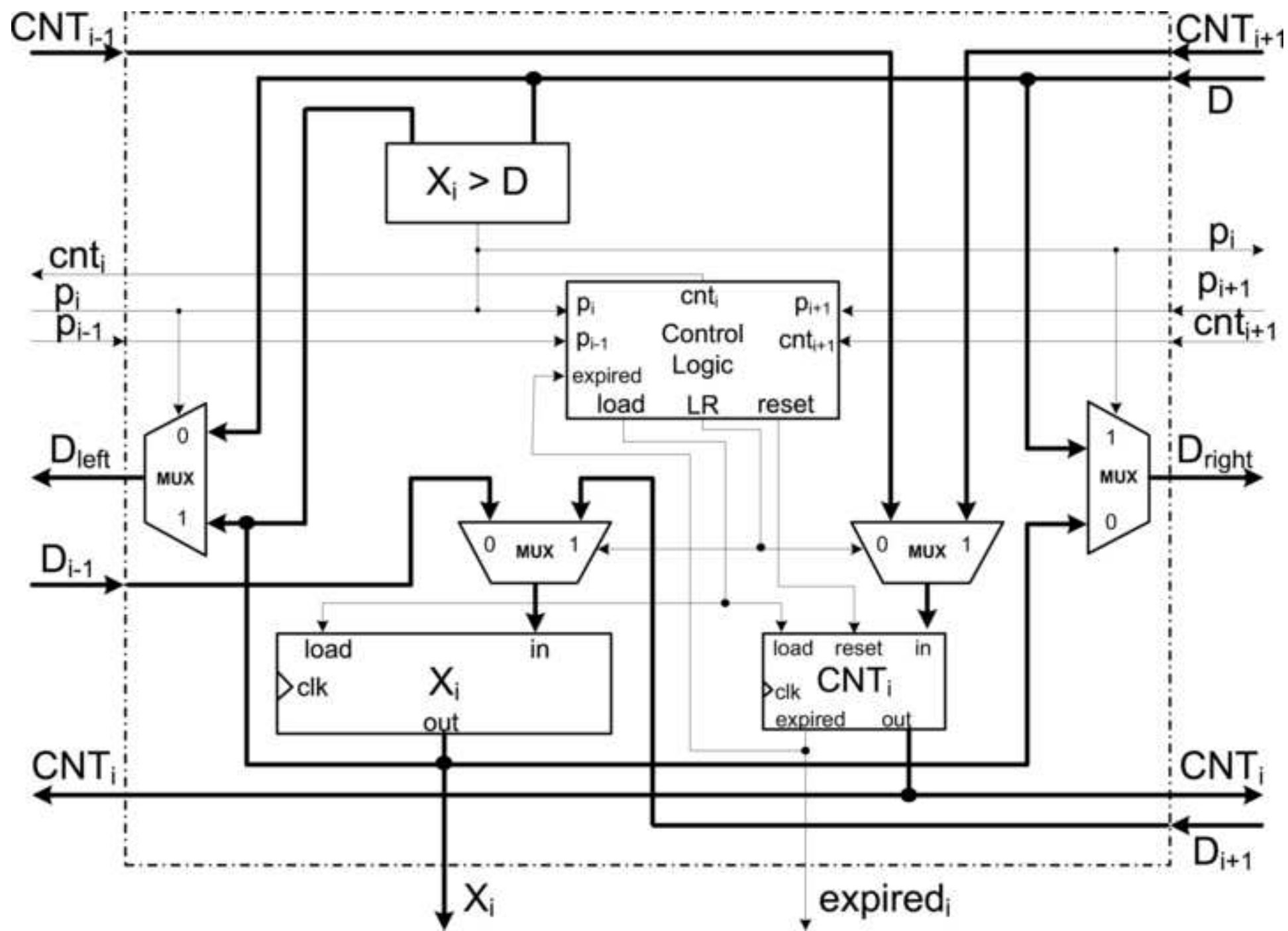


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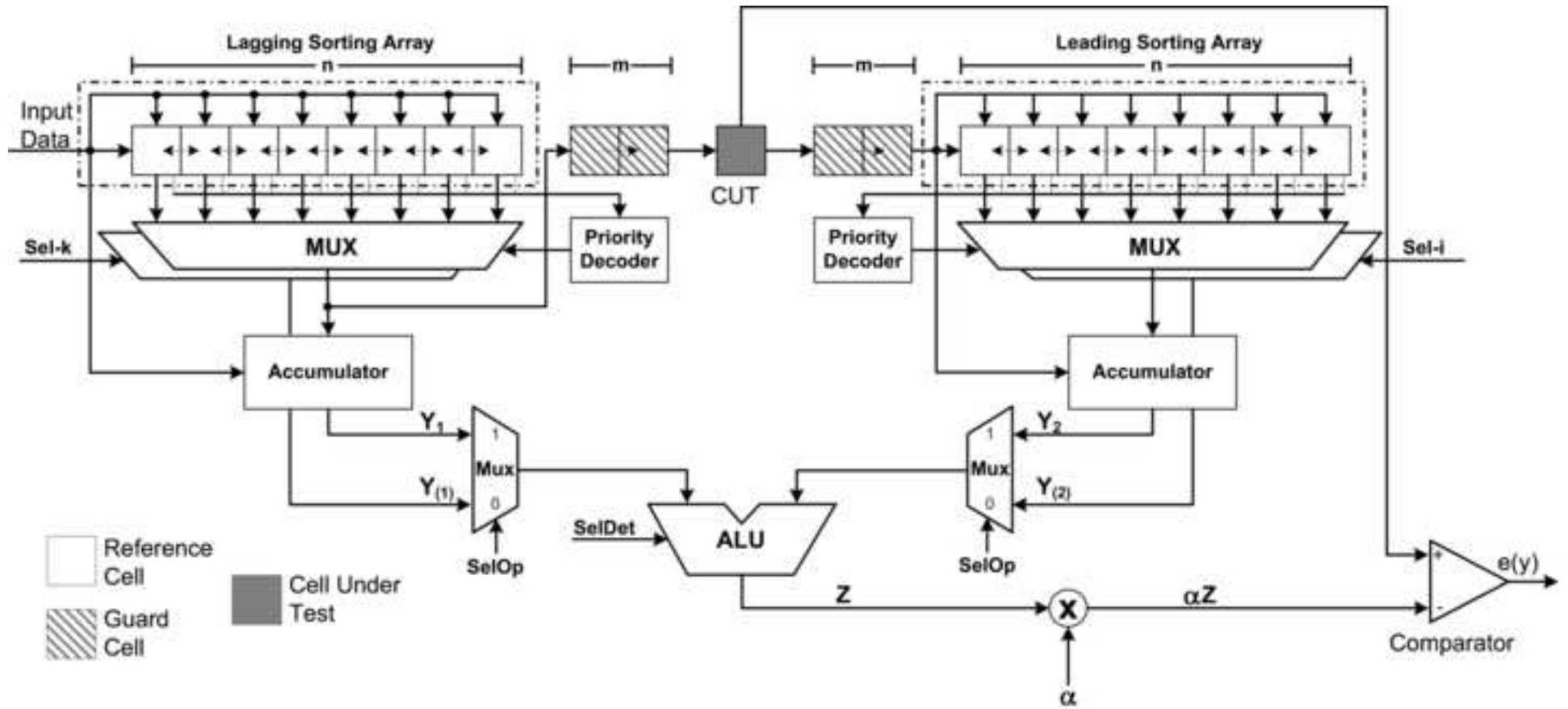


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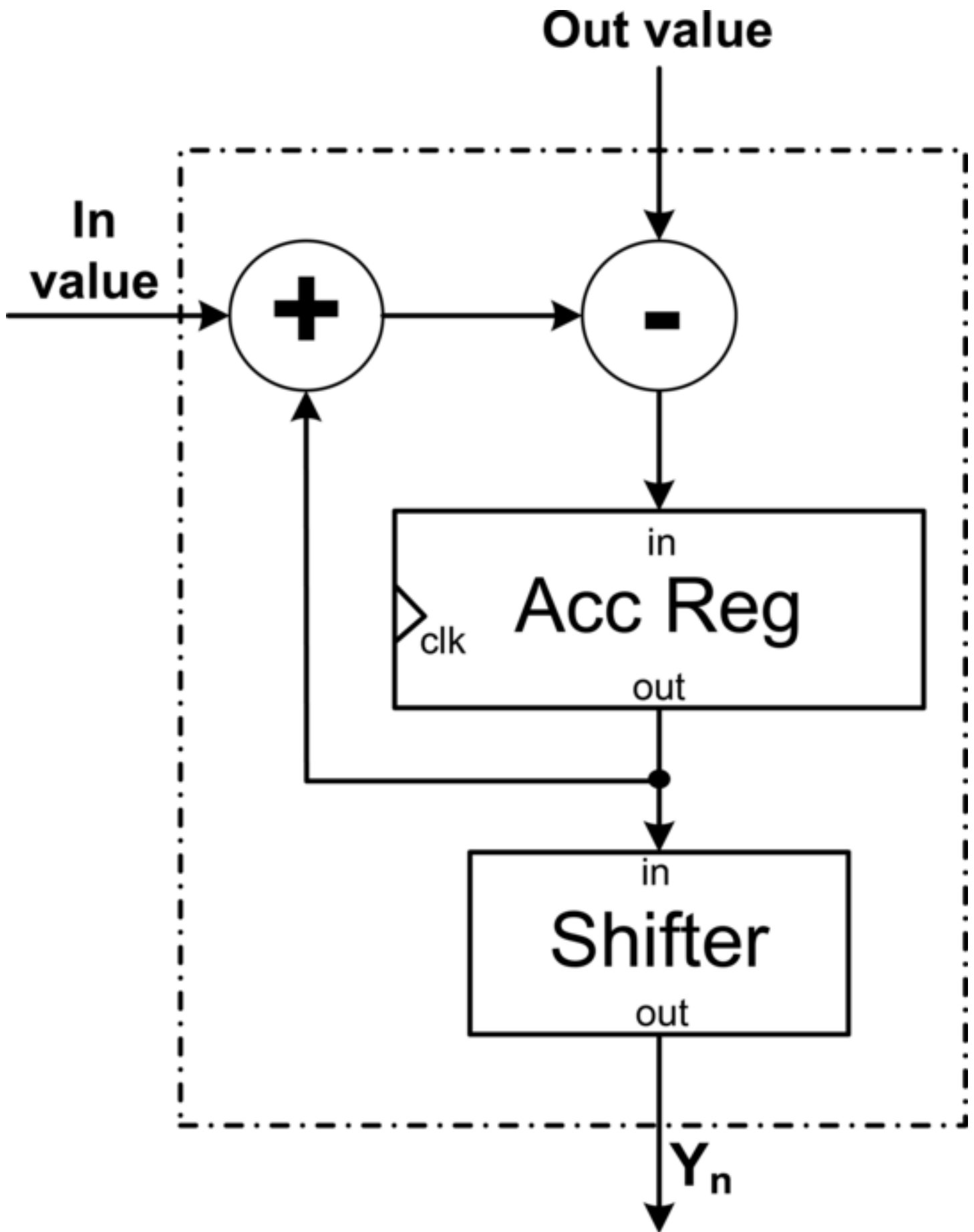


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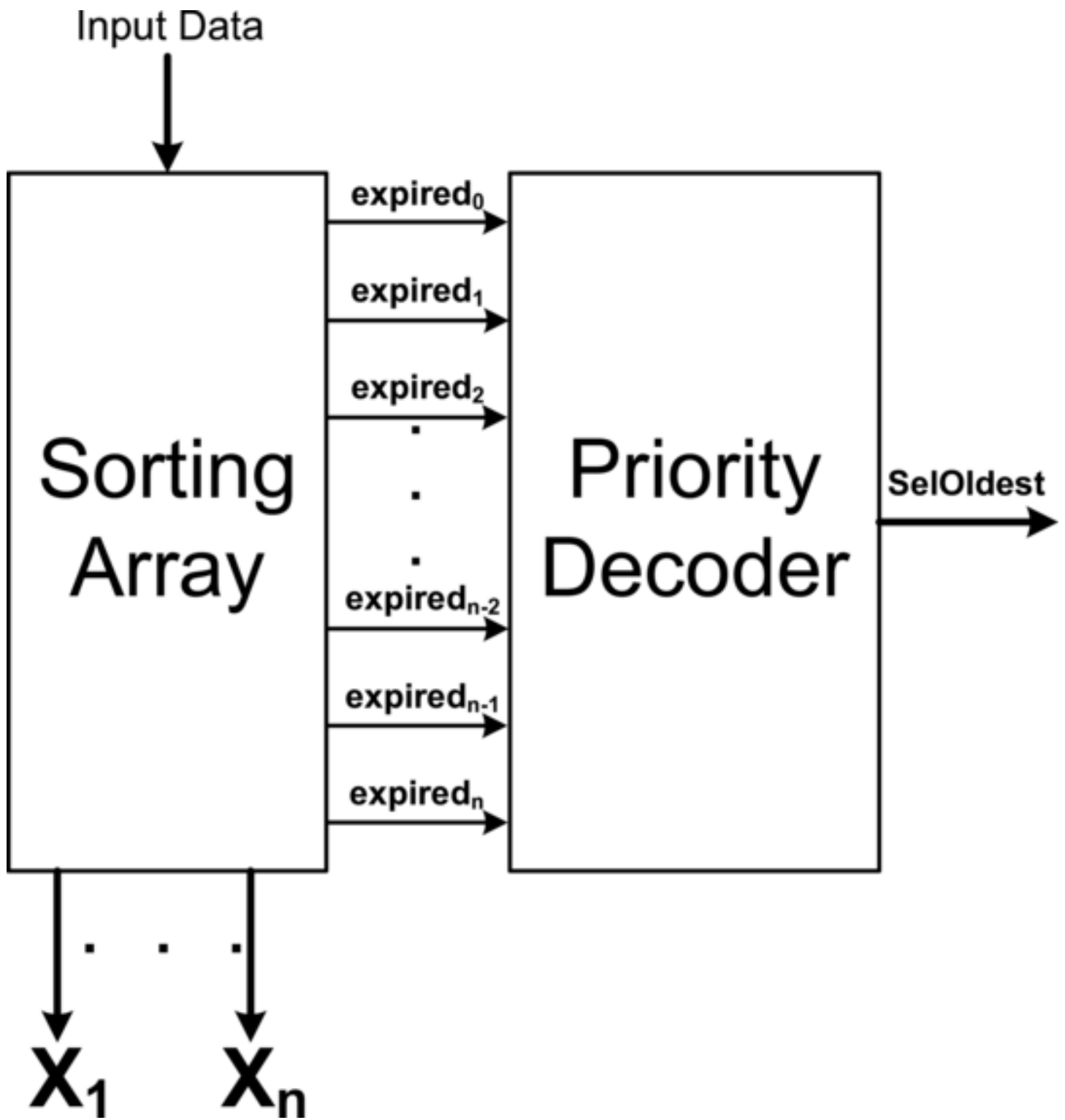


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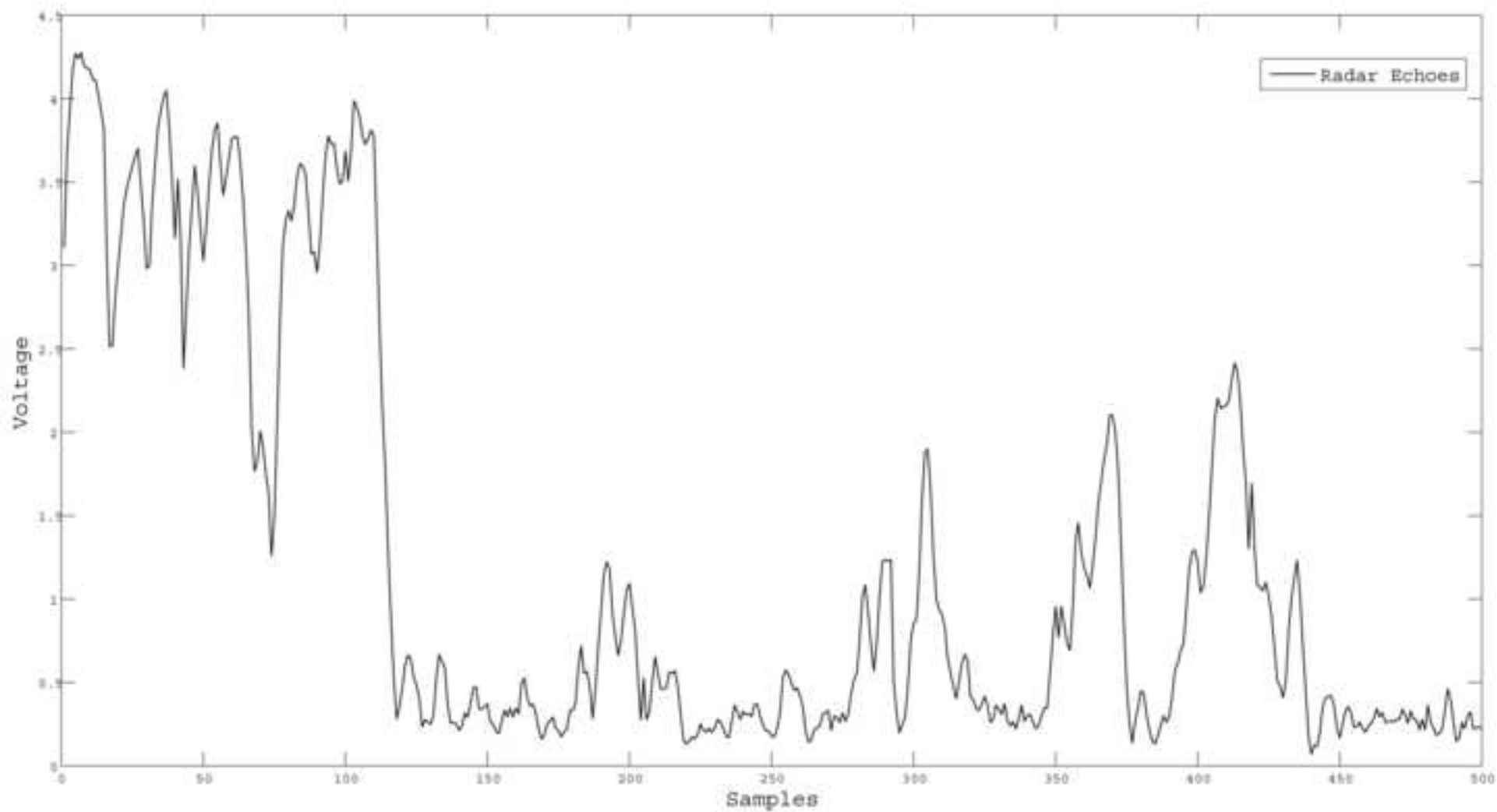




Figure 10a  
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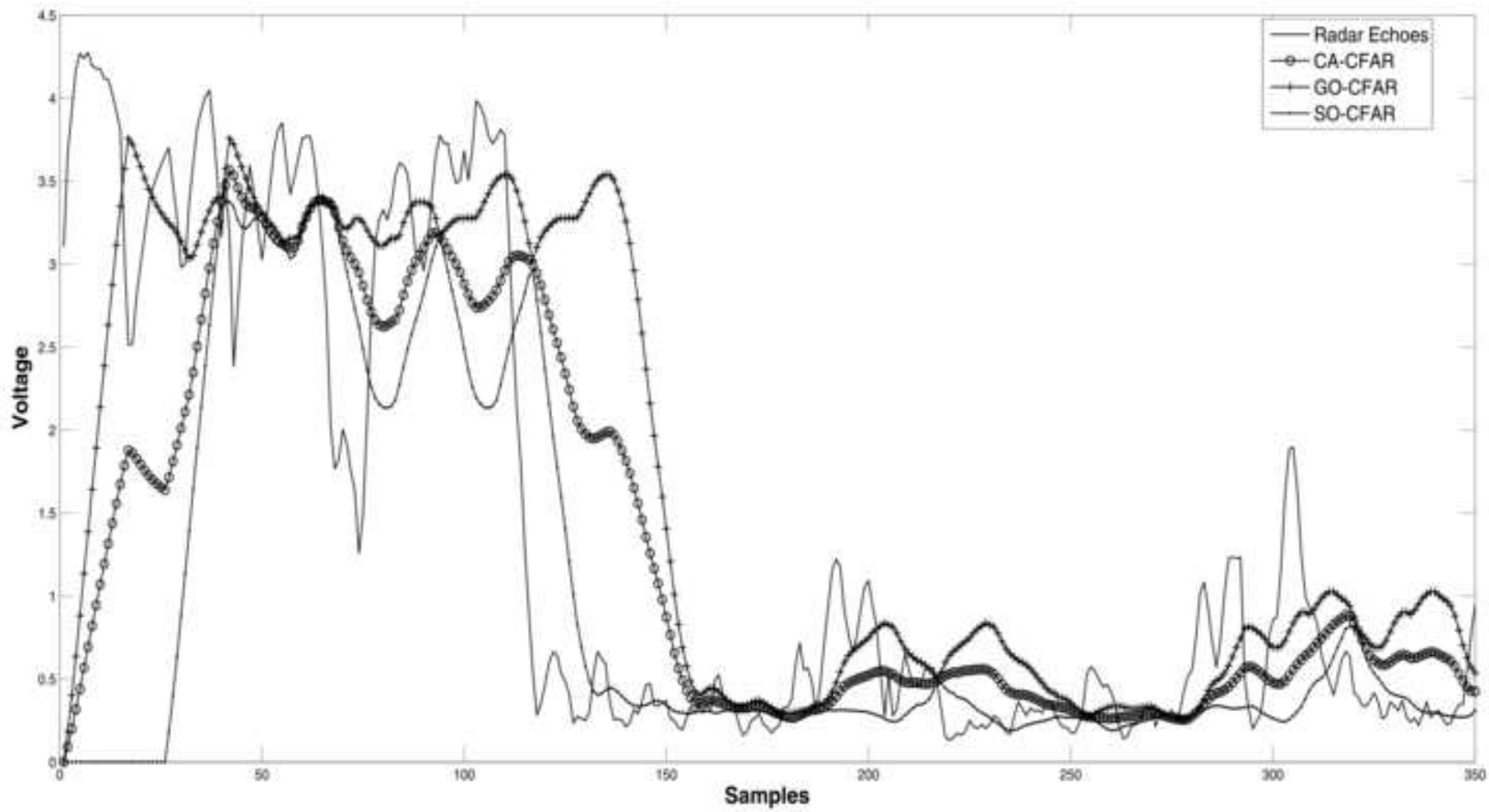


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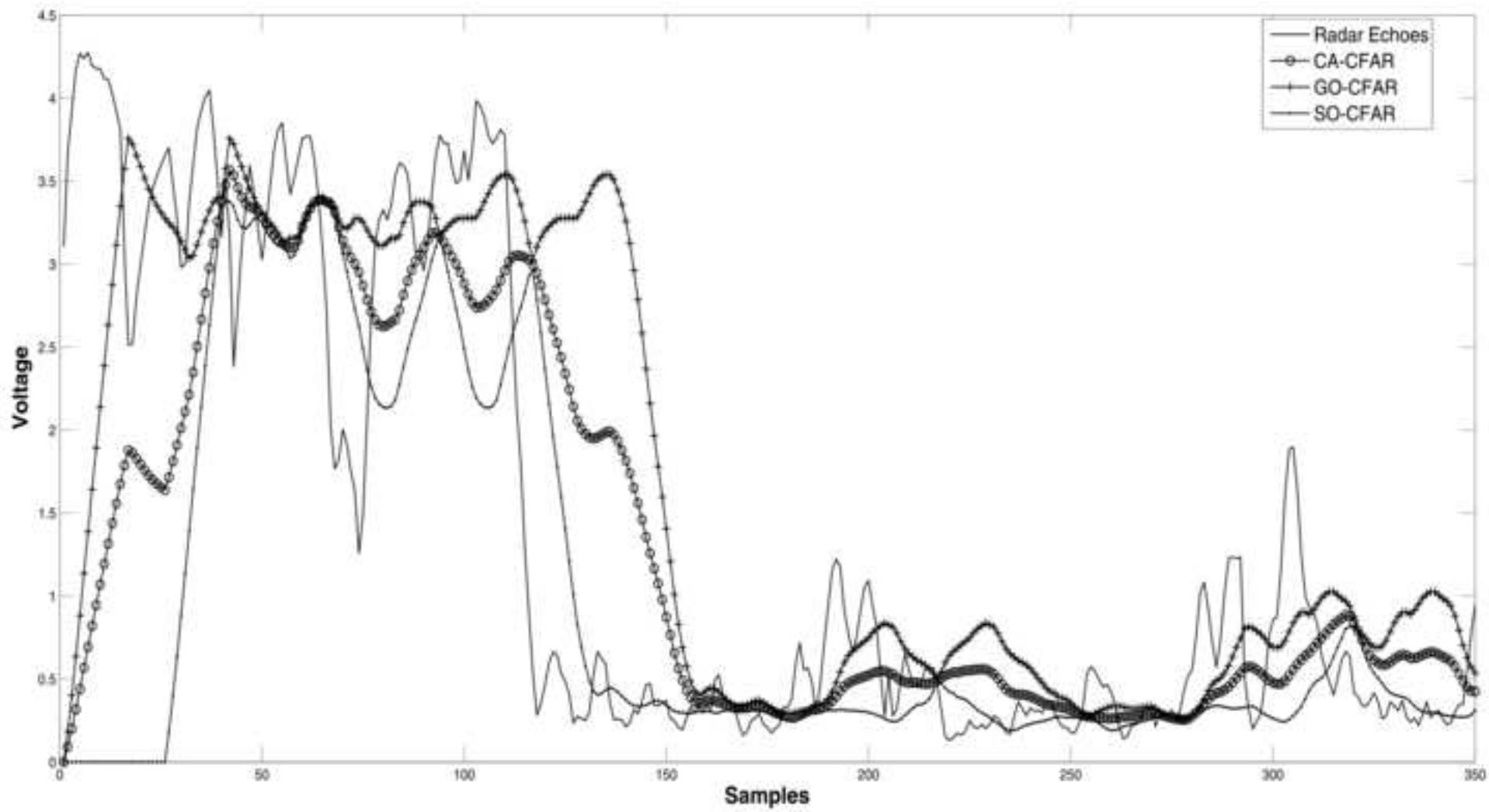


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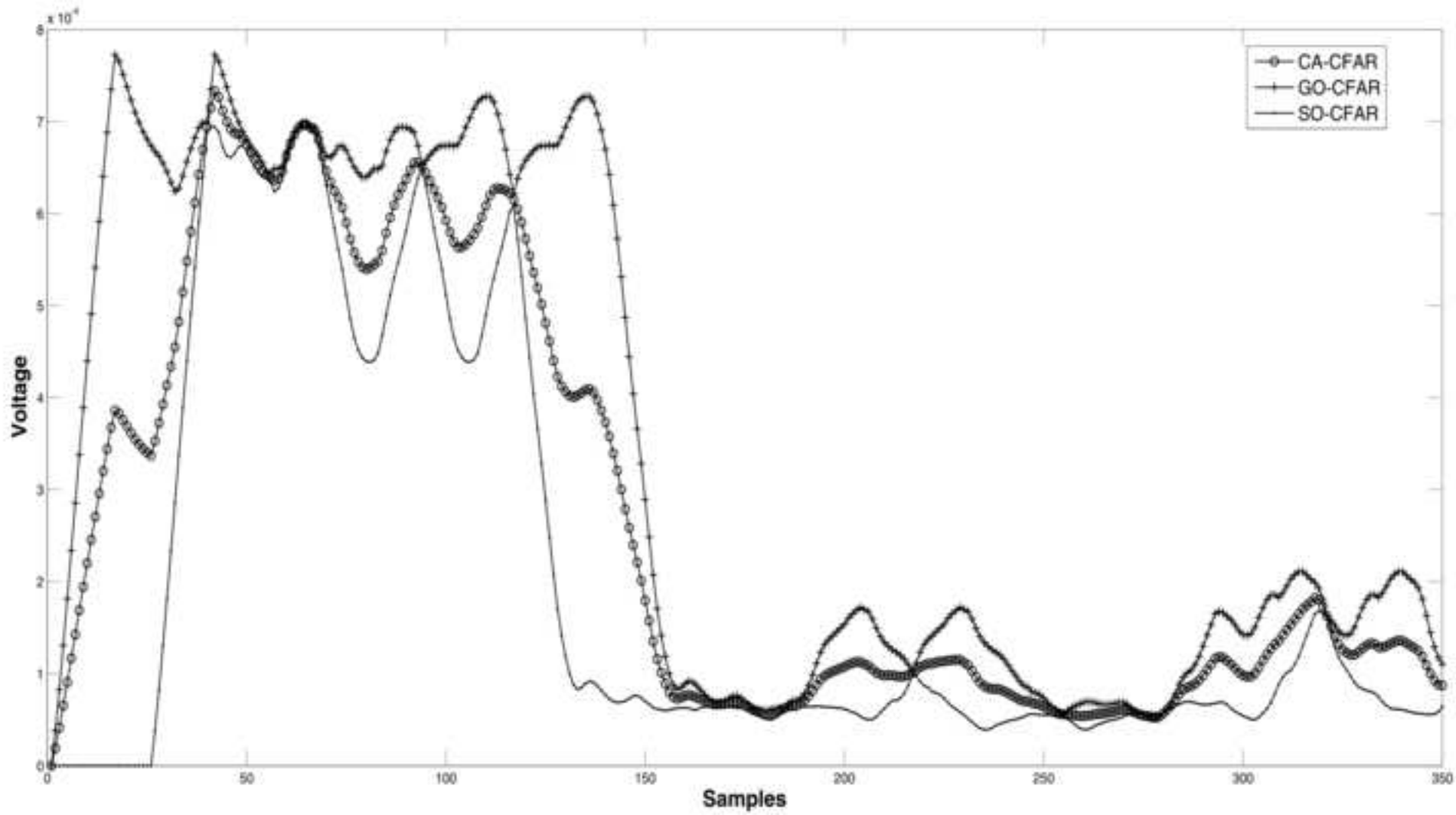


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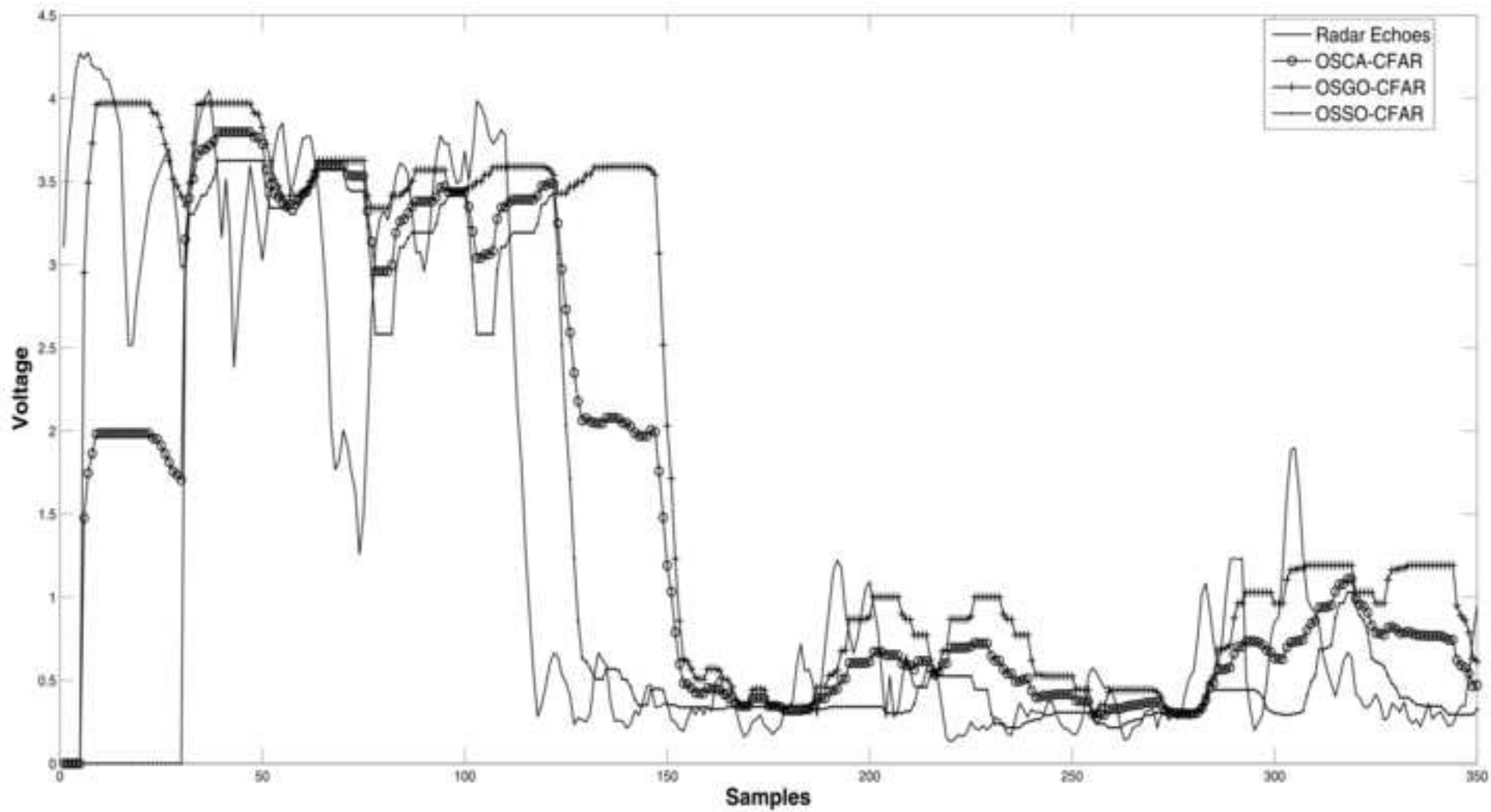


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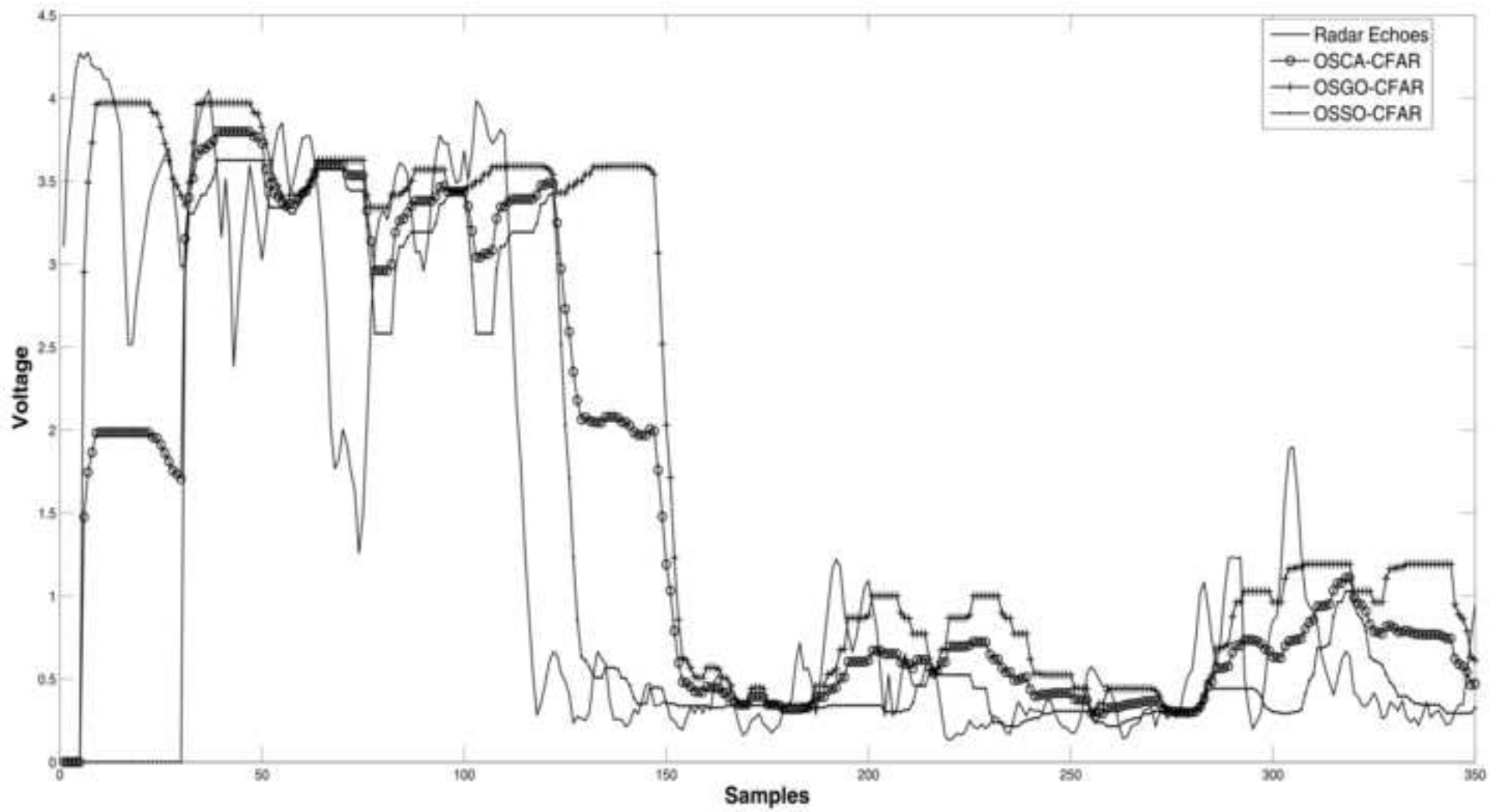


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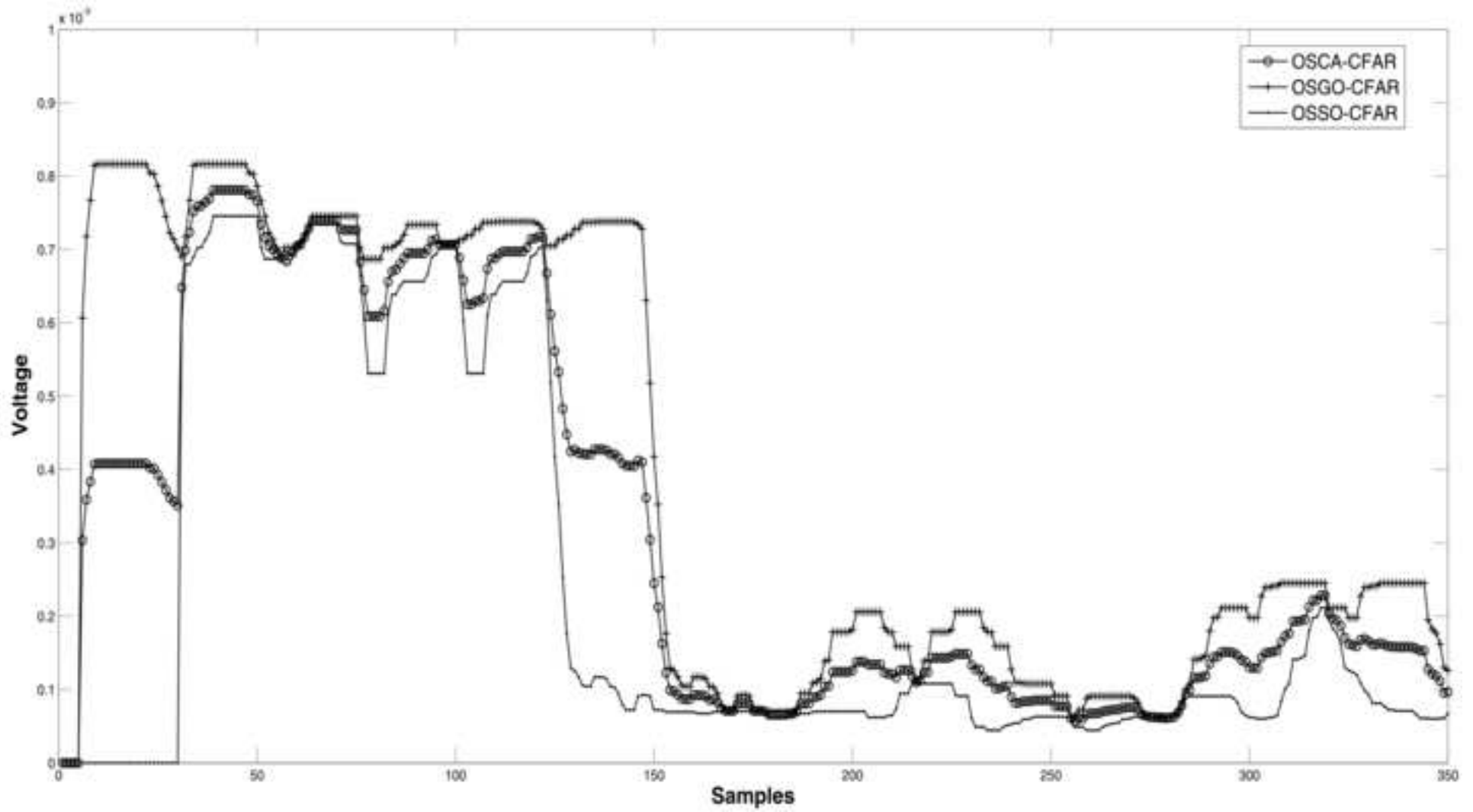


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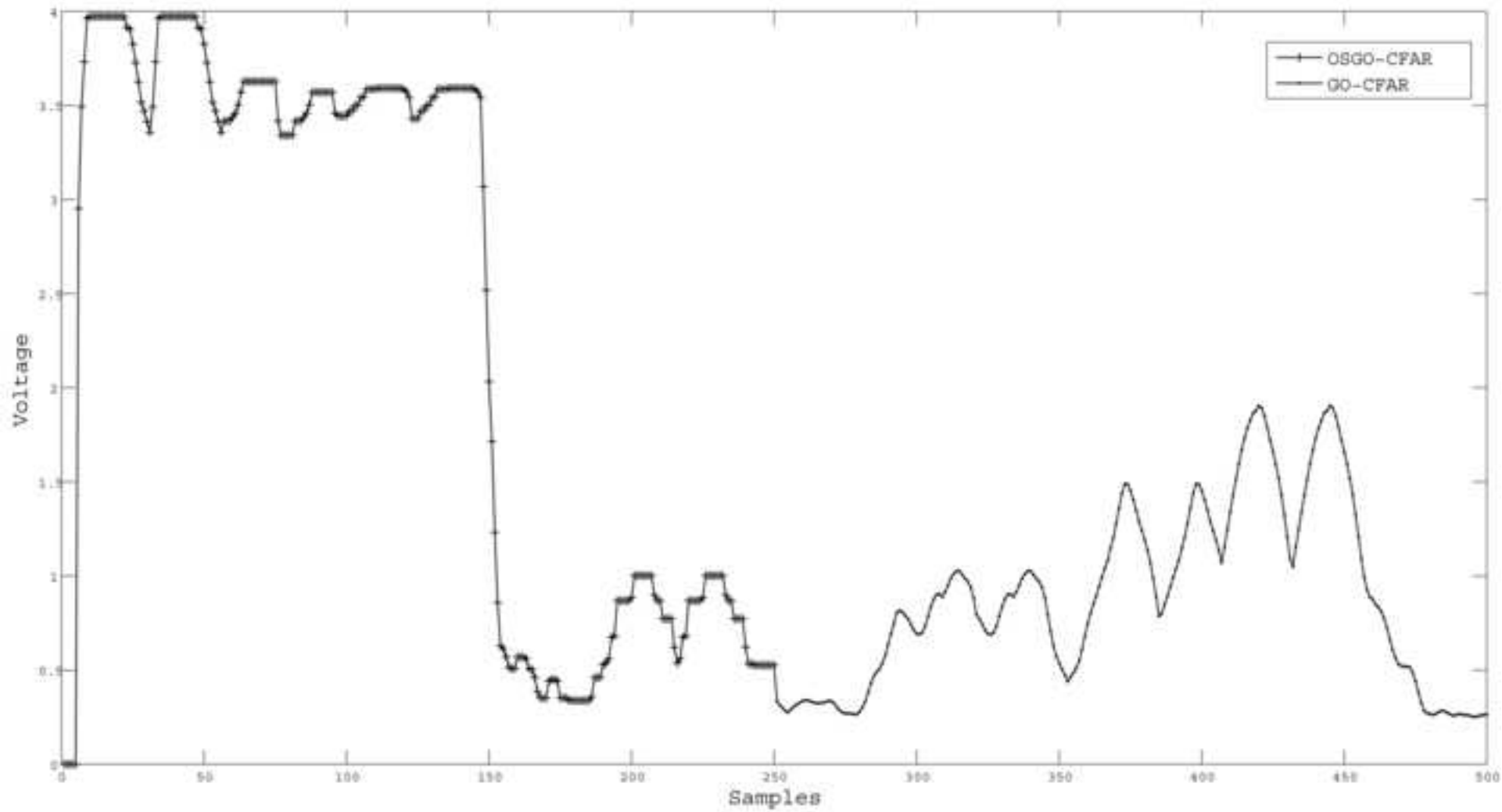


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